

# Ch.5:Backtracking

# 5.1. Backtracking :General method

- Backtracking is a **systematic way** to go through all the possible configurations of a search space.
- In the general case, we assume our **solution** is a vector
$$a = (a[1], a[2], \dots, a[n])$$
where each element  $a[i]$  is selected from a **finite ordered set**  $S[i]$ .

# . Backtracking : General Method

- We build a **partial solution** of length  $k$   
 $a = (a[1], a[2], \dots, a[k])$   
and try to **extend** it by adding another element.
- After extending it, we will **test** whether what we have so far is still possible as a partial solution.

# Backtracking

If it is still a **candidate solution**, great.

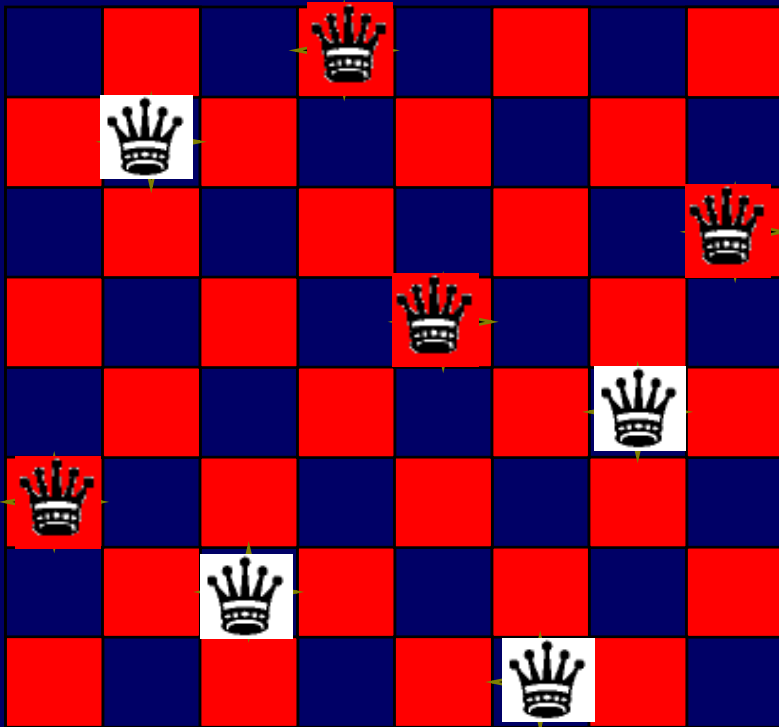
If not, we **delete**  $a[k]$  and **try the next element** from  $S[k]$ .

# Backtracking Concept

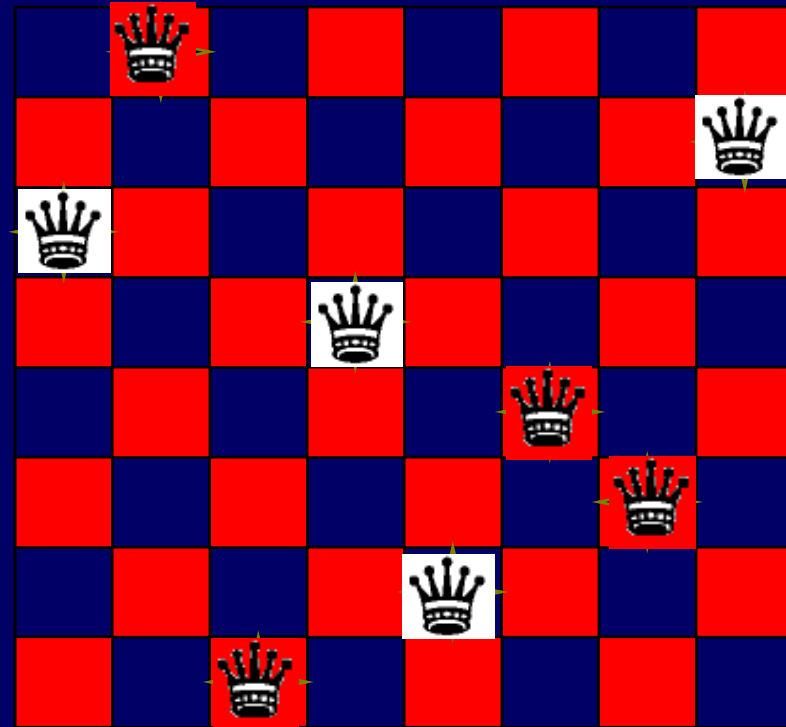
- **Recursion** can be used for elegant and easy implementation of backtracking.
- Backtracking can easily be used to **iterate** through all subsets or permutations of a set.
- Backtracking ensures correctness by enumerating **all possibilities**.
- For backtracking to be efficient, we must **prune** the search space.

# 4.2. Eight Queen Problem (1/7)

Place **8 queens** in a chessboard so that no two queens are in the same row, column, or diagonal.

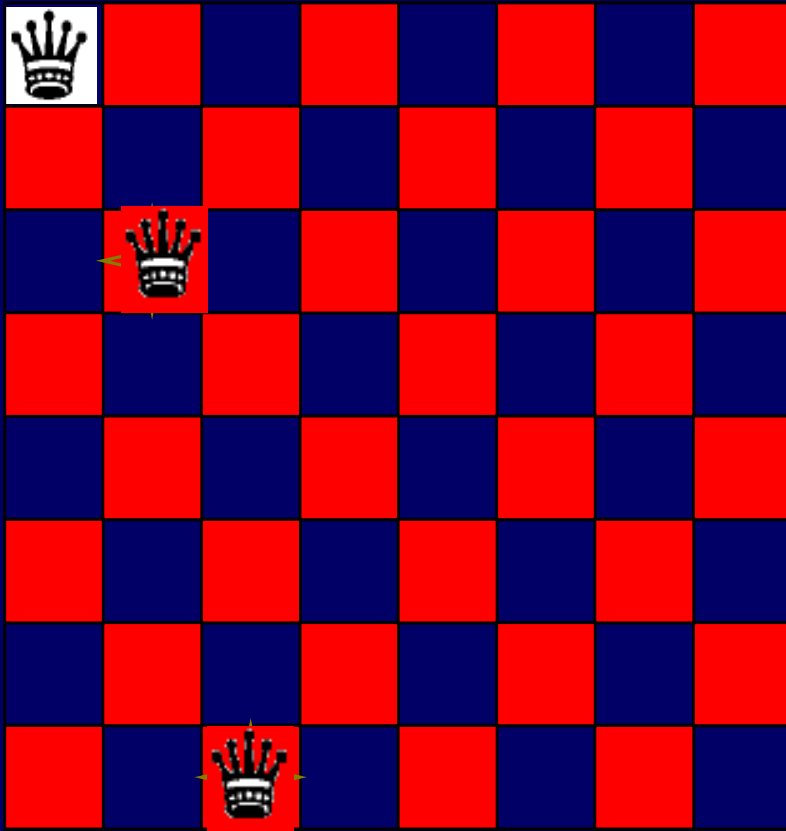


A solution



Not a solution

## 4.2. Eight Queen Problem (2/7)

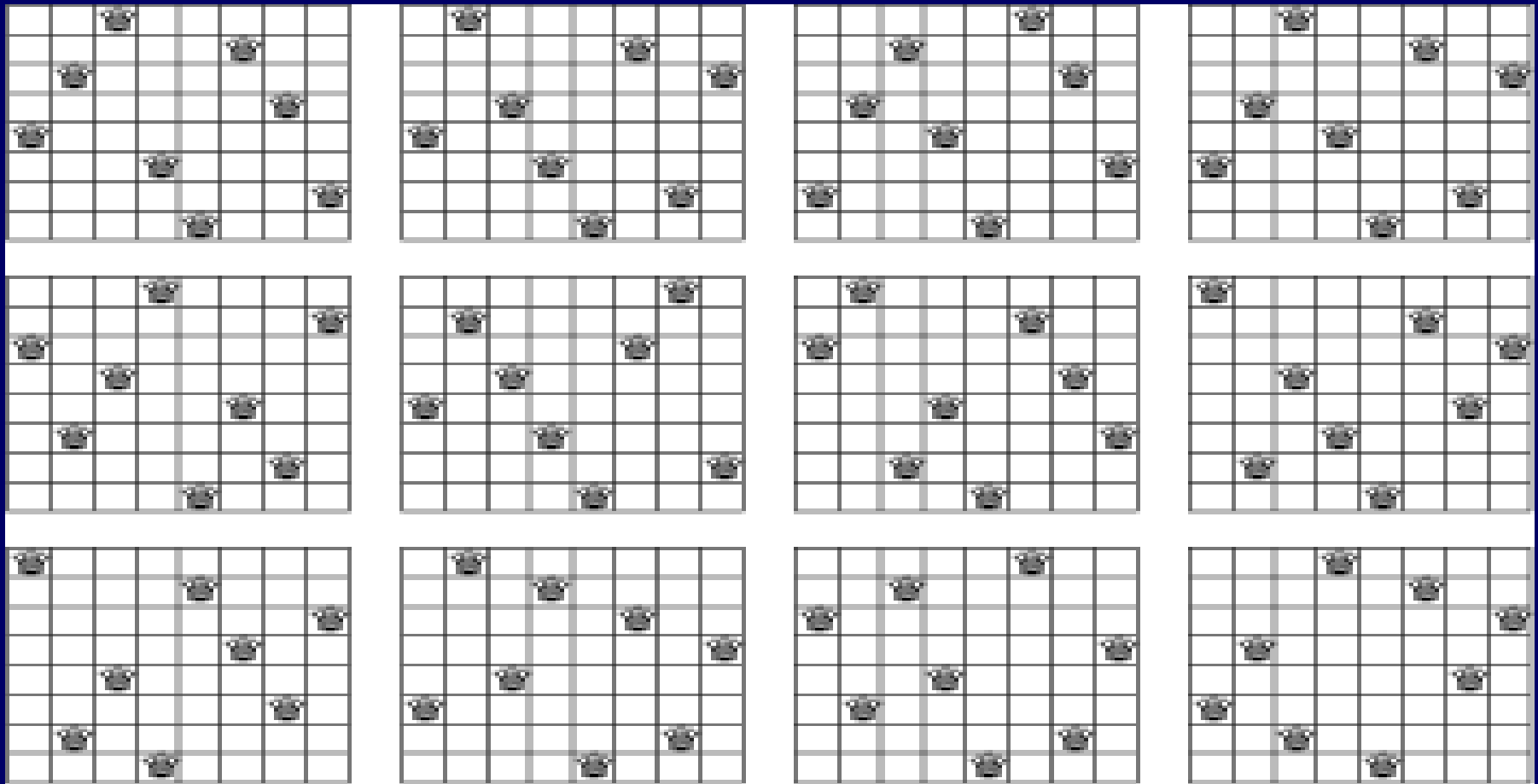


Suppose two queens are placed at position  $(i,j)$  &  $(k,l)$ . Two queens will attack each other if  $i-j=k-l$  or  $i+j=k+l$ . Which is same as  $j-l=i-k$  &  $j-l=k-i$ .

→  $64^8$  states with 8 queens

# 4.2. Eight Queen Problem ()

Some solutions from 92 Solutions





# Can a new queen be placed?

Algorithm Place(k,i)

{

for j=1 to k-1 do

if((x[j]==)or (Abs x[j]-i)=Abs(j-k))

then return false;

return true;

}

# All solution to the n-queen problem

Algorithm NQueen(k, n)

```
{  
  for i= 1 to n do  
  {  
    if Place(k,i) then  
    {  
      x[k]:=i;  
      if(k=n) then write(x[1:n]);  
      else Nqueen(k+1,n);  
    }  
  }  
}
```

# Analysis of 8-Queen problem

If we consider 64 position & reject illegal configuration, no. of configuration will be  $8^6=4,426,165,368$

If we place one queen in one row then  $8^8=16,777,216$

If we reject the position of column, row, diagonal position whose position are guarded then no. of configuration  $8!=40320$

# Application & scope of research

- To develop such an algorithm for eight queen problem whose complexity is less than  $8!$

# Assignment

Q.1) Explain N queen problem.

Q.2) What is attacking position of two queen?

Q.3) What is efficiency of 8-Queen problem?